

Enzyme: High-Performance Automatic Differentiation



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Differentiation Is Key To Machine Learning And Science

- Computing derivatives is key to many algorithms
 - Machine learning (back-propagation, Bayesian inference, uncertainty quantification)
 - Scientific computing (modeling, simulation)
- When working with large codebases or dynamically-generated programs, manually writing derivative functions becomes intractable
- Community has developed tools to create derivatives automatically



Existing AD Approaches

- Differentiable DSL (TensorFlow, PyTorch, DiffTaichi)
 - Provide a new language designed to be differentiated
 - Requires rewriting everything in the DSL and the DSL must support all operations in original code
 - Fast if DSL matches original code well
- Operator overloading (Adept, JAX)
 - Provide differentiable versions of existing language constructs (double => adouble, np.sum => jax.sum)
 - May require writing to use non-standard utilities
 - · Often dynamic: storing instructions/values to later be interpreted

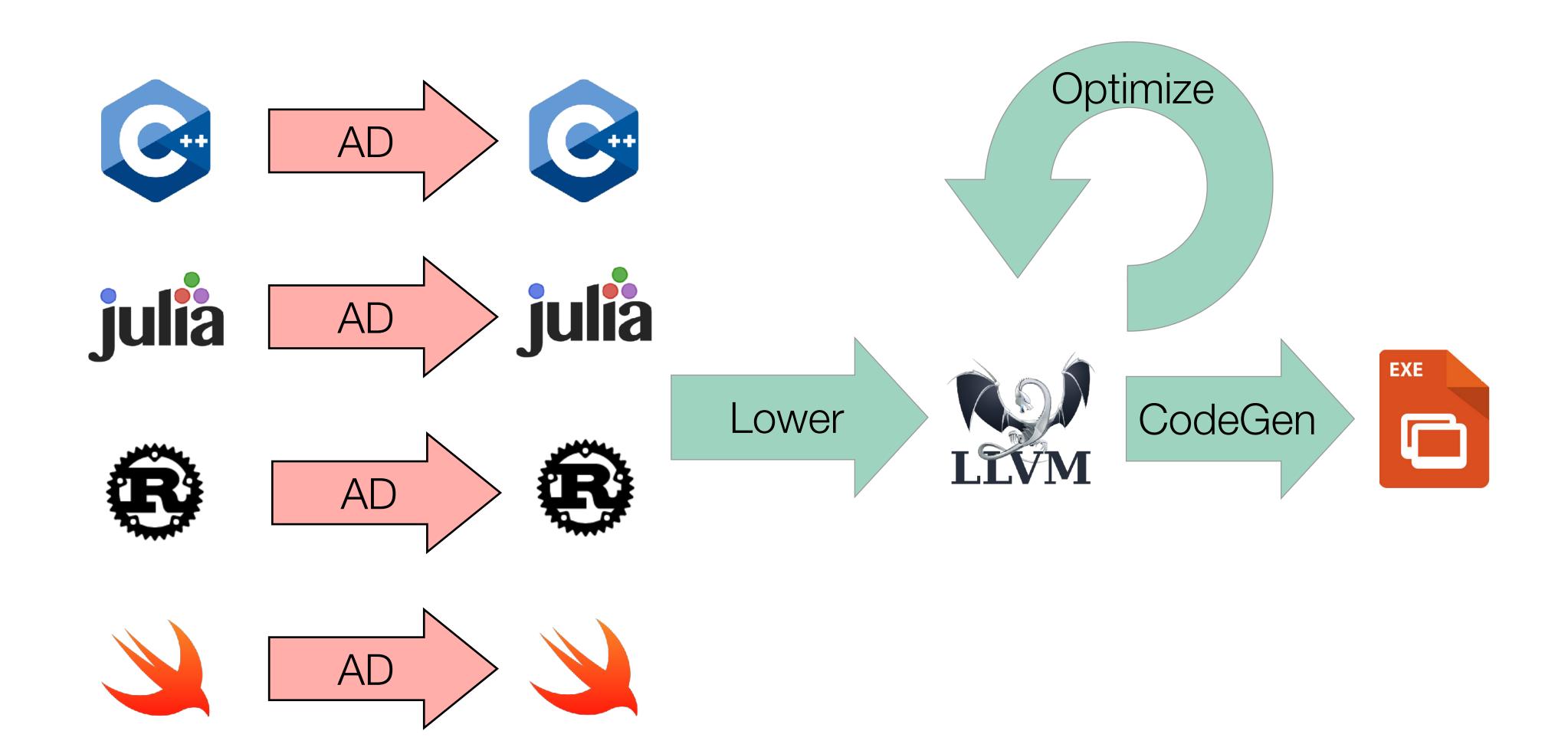


Existing AD Approaches

- Source rewriting
 - Statically analyze program to produce a new gradient function in the source language
 - Re-implement parsing and semantics of given language
 - Requires all code to be available ahead of time
 - Difficult to use with external libraries



Existing Automatic Differentiation Pipelines





Case Study: Vector Normalization

```
//Compute magnitude in O(n)
double mag(double[] x);

//Compute norm in O(n^2)
void norm(double[] out, double[] in) {
  for (int i=0; i<n; i++) {
    out[i] = in[i] / mag(in);
  }
}</pre>
```

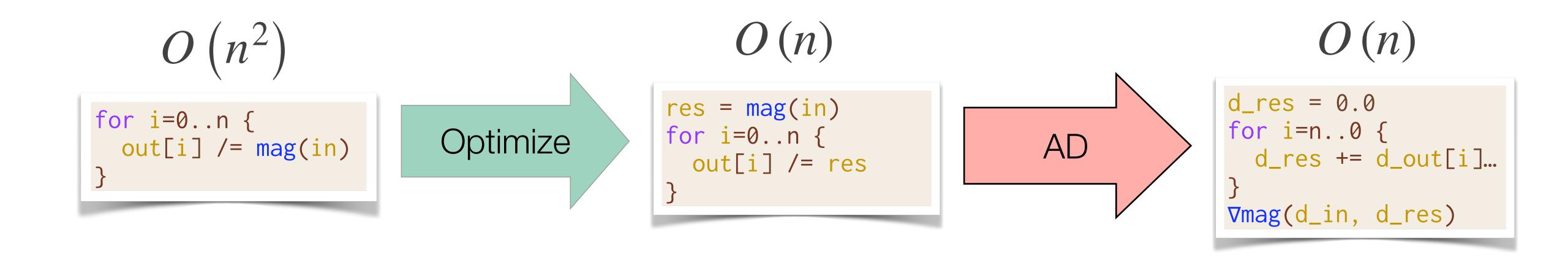


Case Study: Vector Normalization

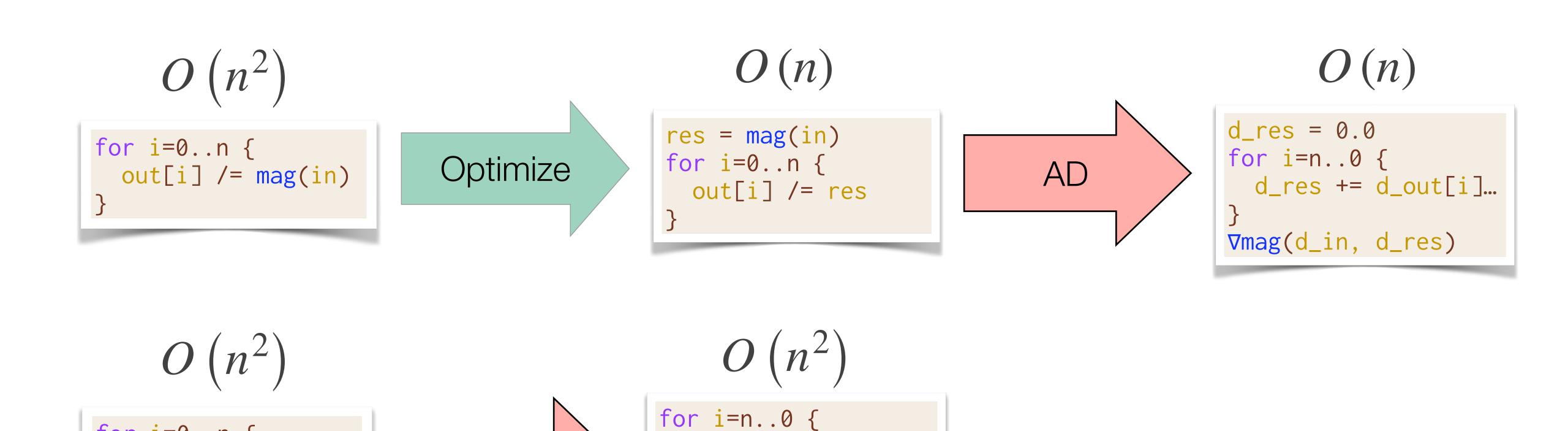
```
//Compute magnitude in O(n)
double mag(double[] x);

//Compute norm in O(n)
void norm(double[] out, double[] in) {
  double res = mag(in);
  for (int i=0; i<n; i++) {
    out[i] = in[i] / res;
  }
}</pre>
```





AD

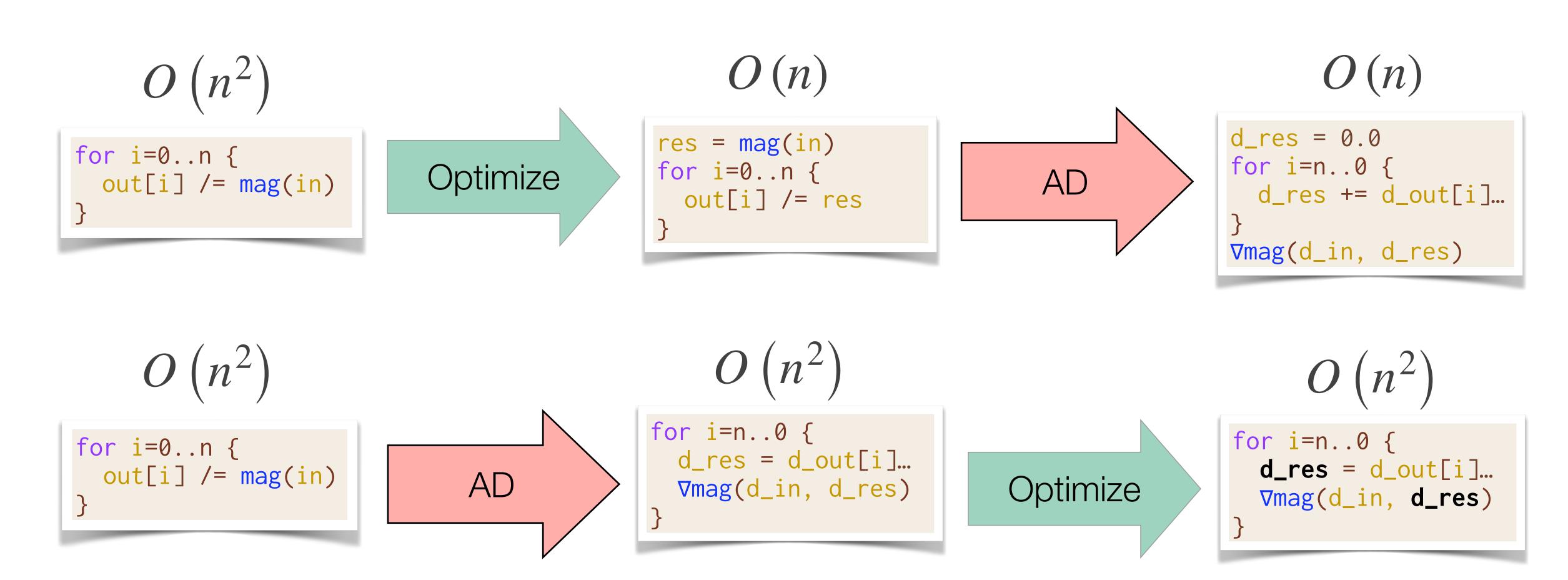


d_res = d_out[i]...

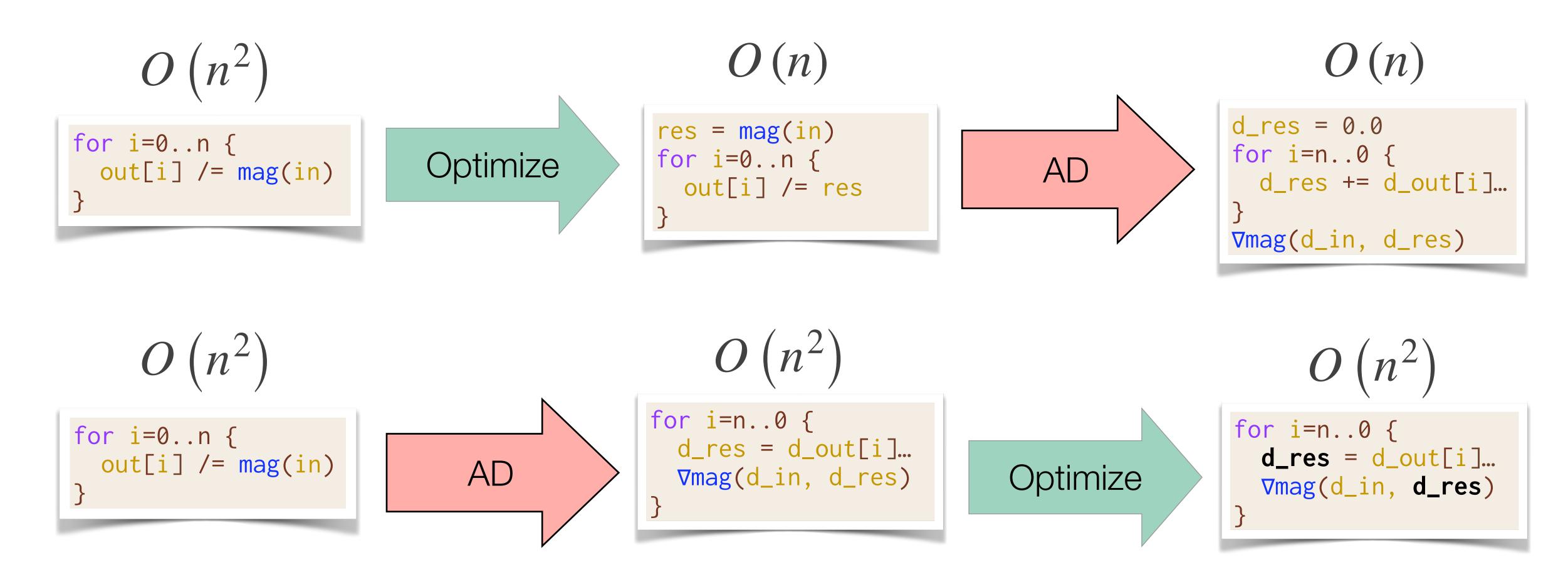
∇mag(d_in, d_res)

for i=0..n {

out[i] /= mag(in)

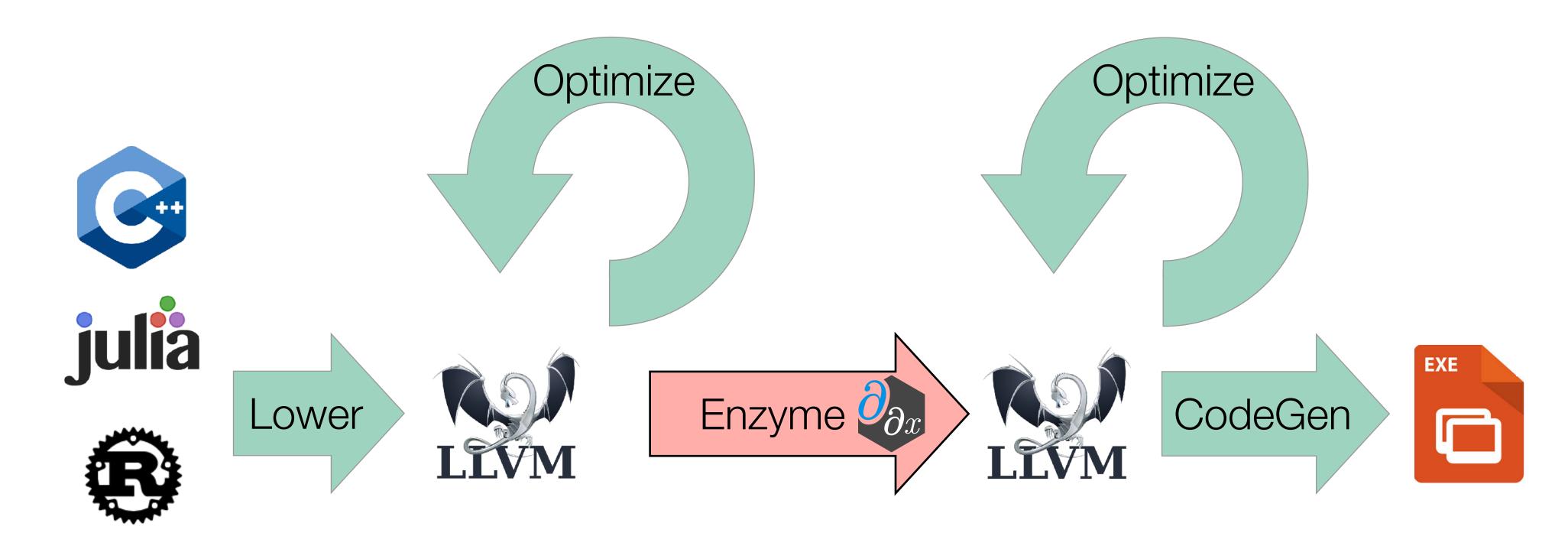


Differentiating after optimization can create asymptotically faster gradients!





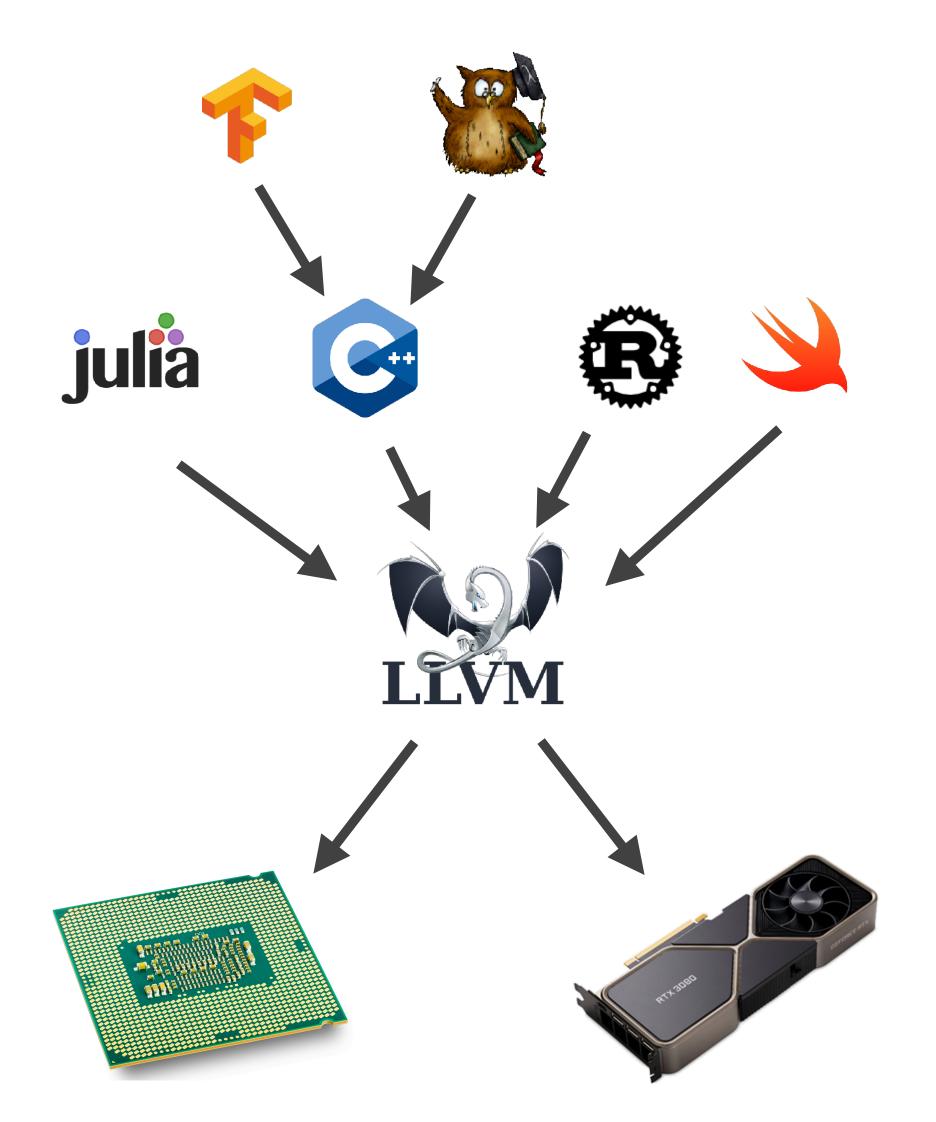
Performing AD at low-level lets us work on optimized code!





Why Does Enzyme Use LLVM?

- Generic low-level compiler infrastructure with many frontends
 - "Cross platform assembly"
 - · Many backends (CPU, CUDA, AMDGPU, etc)
- Well-defined semantics
- Large collection of optimizations and analyses





Case Study: ReLU3

C Source

```
double relu3(double x) {
  double result;
  if (x > 0)
    result = pow(x, 3);
  else
    result = 0;
  return result;
}
```

Enzyme Usage

```
double diffe_relu3(double x) {
  return __enzyme_autodiff(relu3, x);
}
```

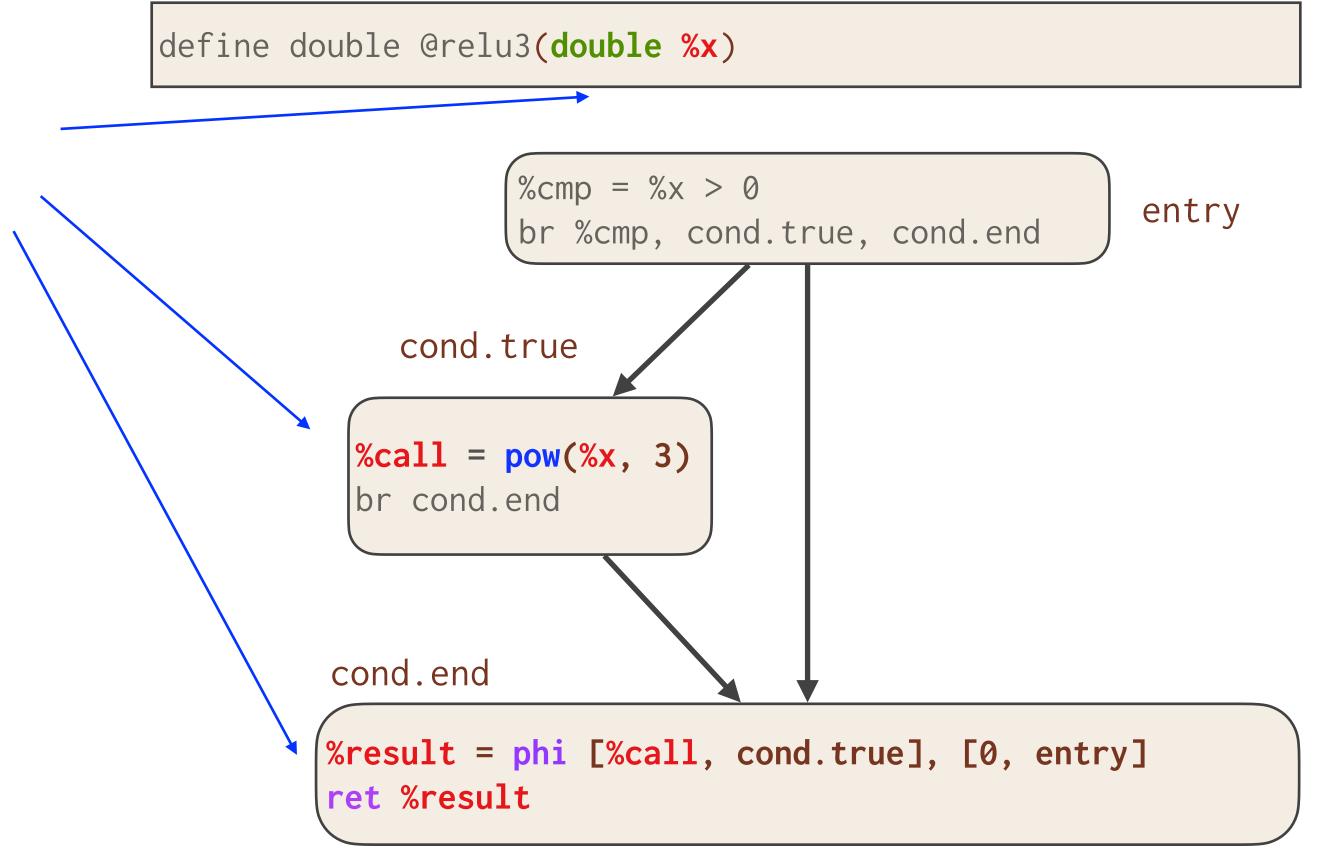
LLVM

```
define double @relu3(double %x)
               entry
                      %cmp = %x > 0
                      br %cmp, cond.true, cond.end
cond.true
    %call = pow(%x, 3)
    br cond.end
           %result = phi [%call, cond.true], [0, entry]
cond.end
            ret %result
```



Case Study: ReLU3

Active Instructions





```
define double @diffe_relu3(double %x, double %differet)
                                                           Allocate & zero
                 alloca %result' = 0.0
                 alloca %call'
                               = 0.0
         entry
                                                       shadow memory for
                 alloca %x'
                               = 0.0
                 %cmp = %x > 0
                                                             active values
                 br %cmp, cond.true, cond.end
cond.true
                                                                             cond.end
  %call = pow(%x, 3)
                            %result = phi [%call, cond.true], [0, entry]
  br cond.end
                             ; deleted return
                            %result' = 1.0
                            br reverse_cond.end
```



```
define double @diffe_relu3(double %x, double %differet)
                  alloca %result' = 0.0
                                                       Compute adjoints
                  alloca %call' = 0.0
         entry
                  alloca %x' = 0.0
                                                    for active instructions
                  %cmp = %x > 0
                  br %cmp, cond.true, cond.end
cond.true
                                                                                 cond.end
     %call = pow(%x, 3)
                              %result = phi [%call, cond.true], [0, entry]
     br cond.end
                                deleted return
                              %result' = 1.0
                              br reverse_cond.end
 reverse_cond.true
                                 %tmp_res' = load %result'
                                 %call' += if %x > 0 then %tmp_res' else 0
                                                                              reverse_cond.end
% df = 3 * pow(%x, 2)
                                  store %result' = 0.0
%tmp_call' = load %call
                                  br %cmp, reverse_cond.true, reverse_entry
%x' += %df * %tmp_call'
store %call' = 0.0
br reverse_entry
                                         \%0 = load \%x'
                                                                 reverse_entry
                                         ret %0
```



```
define double @diffe_relu3(double %x, double %differet)
                  alloca %result' = 0.0
                                                         Compute adjoints
                  alloca %call' = 0.0
         entry
                  alloca %x' = 0.0
                                                       for active instructions
                  %cmp = %x > 0
                  br %cmp, cond.true, cond.end
cond.true
                                                                                 cond.end
     %call = pow(%x, 3)
                              %result = phi [%call, cond.true], [0, entry]
     br cond.end
                              ; deleted return
                              %result' = 1.0
                              br reverse_cond.end
 reverse_cond.true
                                  %tmp_res' = load %result'
                                 %call' += if %x > 0 then %tmp_res' else 0
                                                                              reverse_cond.end
%df = 3 * pow(%x, 2)
                                  store %result' = 0.0
%tmp_call' = load %call
                                  br %cmp, reverse_cond.true, reverse_entry
%x' += %df * %tmp_call'
store %call' = 0.0
br reverse_entry
                                         \%0 = load \%x
                                                                 reverse_entry
                                         ret %0
```



```
define double @diffe_relu3(double %x)

Post

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```

Essentially the optimal hand-written gradient!

```
double diffe_relu3(double x) {
  double result;
  if (x > 0)
    result = 3 * pow(x, 2);
  else
    result = 0;
  return result;
}
```



Challenges of Low-Level AD

Low-level code lacks information necessary to compute adjoints

```
void f(void* dst, void* src) {
  memcpy(dst, src, 8);
}
```

Challenges of Low-Level AD

- New interprocedural dataflow analysis that detects the underlying type of data
- Each value has a set of memory offsets: type
- Perform series of fixed-point updates through instructions

```
struct Type {
  double;
  int*;
}
x = Type*;
```

```
x Type

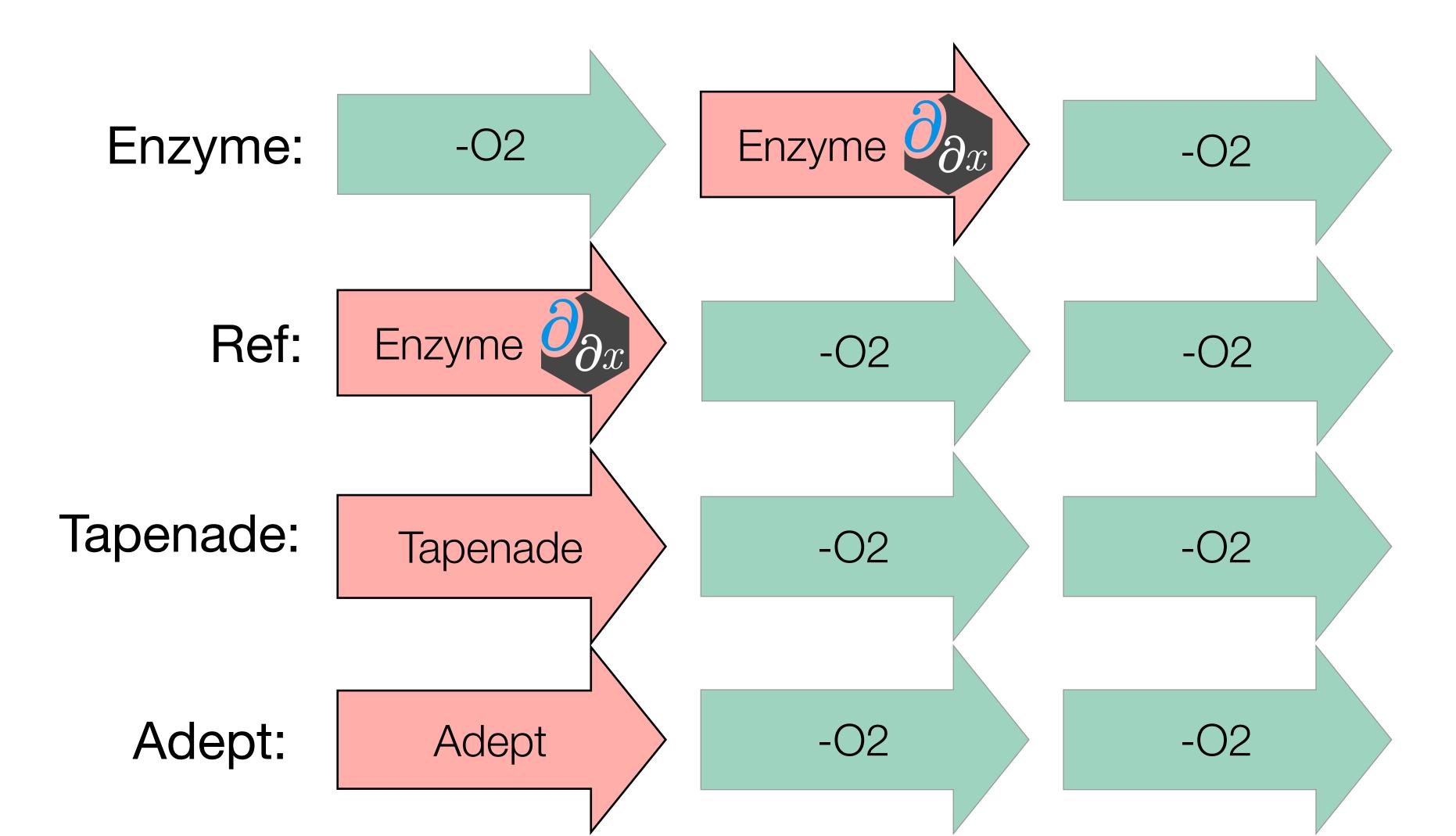
0: Pointer → 0: Double
8: Pointer → 0: Integer
```

```
types(x) = \{[0]: Pointer, [0,0]: Double, [0,8]: Pointer, [0,8,0]: Integer\}
```



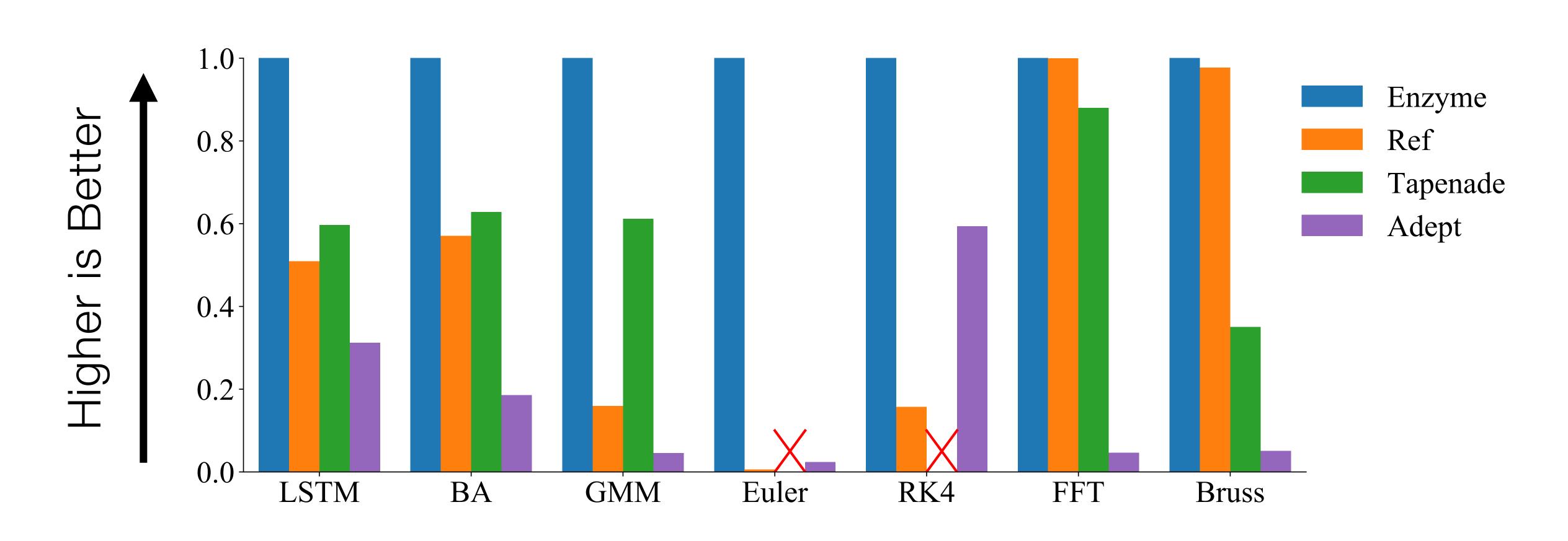
Experimental Setup

Collection of benchmarks from Microsoft's ADBench suite and of technical interest





Speedup of Enzyme



Enzyme is 4.2x faster than Reference!



PyTorch-Enzyme & TensorFlow-Enzyme

```
import torch
from torch_enzyme import enzyme

# Create some initial tensor
inp = ...

# Apply foreign function to tensor
out = enzyme("test.c", "f").apply(inp)

# Derive gradient
out.backward()
print(inp.grad)
```

```
// Input tensor + size, and output tensor
void f(float* inp, size_t n, float* out);

// diffe_dupnoneed specifies not recomputing the output
void diffef(float* inp, float* d_inp, size_t n, float* d_out) {
    __enzyme_autodiff(f, diffe_dup, inp, d_inp, n, diffe_dupnoneed, (float*)0, d_out);
}
```





- Tool for performing reverse-mode AD of statically analyzable LLVM IR
- Differentiates code in a variety of languages (C, C++, Fortran, Julia, Rust, Swift, etc)
- 4.2x speedup over AD before optimization on CPU
- State-of-the art performance with existing tools
- First general purpose reverse-mode GPU AD
- Novel GPU and AD-specific optimizations improve runtime by several orders of magnitude
- · PyTorch-Enzyme & TensorFlow-Enzyme lets researchers use foreign code in ML workflow



Enzyme on the GPU

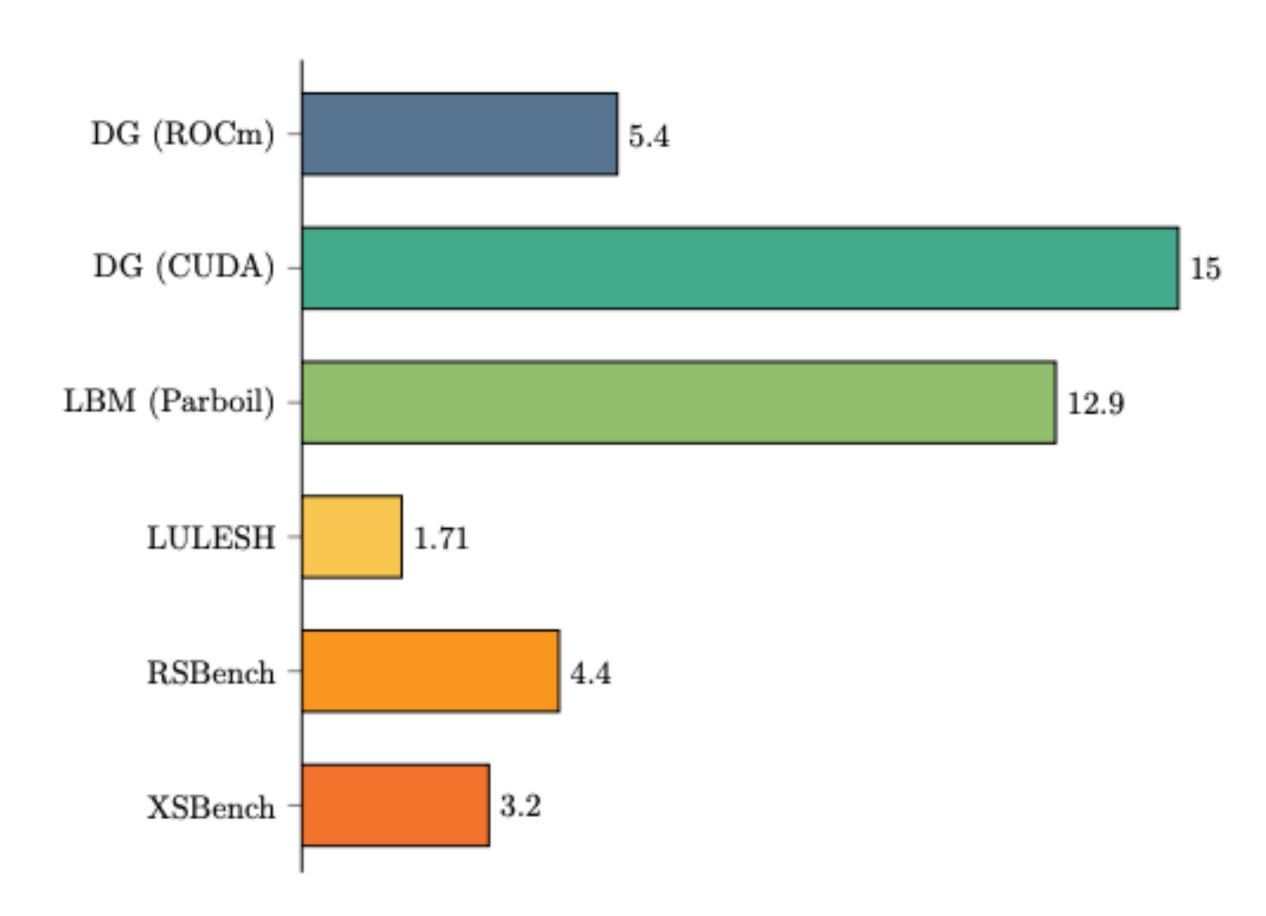
- Care must be taken to both ensure correctness and maintain parallelism.
- GPU programs have much lower memory limits. Performance is highly dependent on the number of memory transfers.
- Without first running optimizations reversemode AD of large kernels is intractable (OOM).
- Novel GPU and AD-specific optimizations can make a difference of several orders of magnitude when computing gradients.

Test	Overhead
Forward	1
AD, Optimized	4.4
AD, No CacheLICM	343.7
AD, Bad Recompute Heuristic	1275.6
AD, No Inlining	6372.2
AD, No PreOptimization	OOM



GPU Gradient Overhead

- Evaluation of both original code and derivative of all inputs (forward or numeric differentiation requires 1 evaluation per input):
 - DG: discontinuous-galerkin (DG) volume integral (Julia)
 - LBM: particle-based fluid dynamics simulation
 - LULESH: unstructured explicit shock hydrodynamics solver
 - XSBench & RSBench: Monte Carlo simulations of particle transport algorithms (memory & compute bound, respectively)





CUDA Automatic Differentiation

- Enzyme enables differentiation of CPU programs without rewriting them in a DSL.
- Similarly, GPU programs cannot currently be differentiated without being rewritten in a differentiable language (e.g. PyTorch).
- Enzyme enables reverse-mode AD of general existing GPU programs by:
 - Resolving potential data race issues
 - Differentiating parallel control (syncthreads)
 - Differentiating CUDA intrinsics (e.g. threadIdx.x /Ilvm.nvvm.read.ptx.sreg.tid.x)
 - Handling shared memory



Challenges of Parallel AD

Benign read race in forward pass => Write race in reverse pass (undefined behavior)



Parallel Memory Detection

Thread-local memory

Same memory location across all threads

Others [always legal fallback]

Non-atomic load/store

Parallel Reduction

Atomic increment

```
__device__
void f(...) {

  // Thread-local var
  double y;

...

  d_y += val;
}
```

```
// Same var for all threads
double y;

__device__
void f(...) {
    ...
    reduce_add(&d_y, val);
}
```

```
__device__
// Unknown thread-aliasing
void f(double* y) {
    ...
    atomic { d_y += val; }
}
```

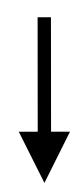


CUDA Automatic Differentiation

```
%res = load %ptr
```

```
%tmp = load %d_res
store %d_res = 0
atomic %d_ptr += %tmp
```

```
store %ptr = %val
```



```
%tmp = load %d_ptr
store %d_ptr = 0
load/store %d_val += %tmp
```

- Shadow Registers %d_res and %d_val are *thread-local* as they shadow thread-local registers.
 - No risk of races and no special handling required.
- Both %ptr and shadow %d_ptr might be raced upon and require analysis.



Differentiation of SyncThreads

- Sync is only necessary if A and B may write to the same memory
- Four cases for what sync could represent:
 - 1. All stores in A must complete prior to a load in B
 - 2. All loads in A must complete prior to a store in B
 - 3. All stores in A must complete prior to a stores in B [clobber]
 - 4. All load in A must complete prior to a load in B [unnecessary sync]

```
codeA();
sync_threads;
codeB();
```



CUDA Performance Improvements

- Enzyme may need to cache values from the forward pass for later use in a reverse pass computation
 - When a value needs caching, Enzyme allocates memory (via malloc inside kernel)
 - Potentially quite slow
 - May overwhelm the amount of GPU heap memory

```
void f(float* in, float* out) {
  float tmp;
  for (int i=0; i<N; i++) {
    tmp = compute(in, i);
    out[i] = tmp * tmp + ...;
  }
    Value tmp is overwritten every
  iteration and must be cached</pre>
```

```
void diffe_f(float* in, float* out) {
    float* tmp_cache = malloc(...);

    for (int i=0; i<N; i++) {
        ...
        tmp_cache[i] = tmp;
    }

    for (int i=N-1; i>=0; i--) {
        ...
        d_tmp[0] = 2 * tmp_cache[0] * d_out[i];
        d_compute(...);
    }

    free(tmp_cache);
}
```



Case 1: Store, Sync, Load

```
codeA(); // store %ptr
sync_threads;
codeB(); // load %ptr
diffe_codeB(); // atomicAdd %d_ptr
sync_threads;
diffe_codeA(); // load %d_ptr
               // store %d_ptr = 0
```



 Load of d_ptr must happen after all atomicAdds have completed



Case 2: Load, Sync, Store

```
codeA(); // load %ptr
sync_threads;
codeB(); // store %ptr
diffe_codeB(); // load %d_ptr
               // store %d_ptr = 0
sync_threads;
diffe_codeA(); // atomicAdd %d_ptr
```



 All of the stores of d_ptr will complete prior to any atomicAdds

No cross-thread race here since that's equivalent to a write race in B



Case 3: Store, Sync, Store

```
codeA(); // store %ptr
sync_threads;
codeB(); // store %ptr
diffe_codeB(); // load %d_ptr
               // store %d_ptr = 0
sync_threads;
diffe_codeA(); // load %d_ptr
               // store %d_ptr = 0
```



 All stores to d_ptr in diffe_B will complete prior to diffe_A, ensuring only the clobbering store has its derivative incremented



CUDA Automatic Differentiation

- Most CUDA intrinsics [e.g. threadIdx.x] are inactive and recomputable and thus are incorporated into Enzyme without any special handling
- Derivative of syncthreads is a syncthreads at the corresponding place in reverse pass
- Shared memory is handled by making a second shared memory allocation to act as the shadow for any potentially active uses



CUDA Example

```
__device__ void inner(float* a, float* x, float* y) {
  y[threadIdx.x] = a[0] * x[threadIdx.x];
}
__device__ void __enzyme_autodiff(void*, ...);

__global__ void daxpy(float* a, float* da, float* x, float* dx, float* y, float* dy) {
  __enzyme_autodiff((void*)inner, a, da, x, dx, y, dy);
}
```

```
__device__ void diffe_inner(float* a, float* da, float* x, float* dx, float* y, float* dy) {
   y[threadIdx.x] = a[0] * x[threadIdx.x];

   float dy = dy[threadIdx.x];
   dy[threadIdx.x] = 0.0f;

   float dx_tmp = a[0] * dy;
   atomic { dx[threadIdx.x] += dx_tmp; }

   float da_tmp = x[threadIdx.x] * dy;
   atomic { da[0] += da_tmp; }
}
```



CUDA Performance Improvements

- Introduce optimizations to reduce the use of memory
 - Alias Analysis to determine legality of recomputing an instruction
 - More aggressive alias analysis properties of syncthreads
 - Don't cache unnecessary values
 - Move cache outside of loops when possible
 - Heap-to-stack [and to register]
 - Don't cache memory itself acting as a cache [such as shared memory]



CUDA Evaluation

	Forward Pass	Gradient No Opt	+ Standard Opts	+ Cache Opts
XSBench- CUDA	1.0s	OOM	20.1s	5.0s
RSBench- CUDA	1.9s	OOM	>540s	7.8s





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- · Differentiates code in a variety of languages (C, C++, Fortran, Julia, Rust, Swift, etc)
- 4.2x speedup over AD before optimization
- State-of-the art performance with existing tools
- Differentiate GPU kernels
- · Open Source (enzyme.mit.edu / github.com/wsmoses/Enzyme)
- PyTorch-Enzyme & TensorFlow-Enzyme imports foreign code in ML workflow

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END



Compiler Analyses Better Optimize AD

- Existing
- Alias analysis results that prove a function does not write to memory, we can prove that additional function calls do not need to be differentiated since they cannot impact the output
- Don't cache equivalent values
- · Statically allocate caches when a loop's bounds can be determined in advance



Decomposing the "Tape"

- Performing AD on a function requires data structures to compute
 - All values necessary to compute adjoints are available [cache]
 - Place to store adjoints [shadow memory]
 - Record instructions [we are static]
- · Creating these directly in LLVM allows us to explicitly specify their behavior for optimization, unlike approaches that call out to a library
- For more details look in paper



Conventional Wisdom: AD Only Feasible at High-Level

- Automatic Differentiation requires high level semantics to produce gradients
- Lack of high-level information can hinder performance of low-level AD
 - "AD is more effective in high-level compiled languages (e.g. Julia, Swift, Rust, Nim) than traditional ones such as C/C++, Fortran and LLVM IR [...]" -Innes^[1]

[1] Michael Innes. Don't Unroll Adjoint: Differentiating SSA-Form Programs. arXiv preprint arXiv:1810.07951, 2018



Differentiation Is Key To Machine Learning

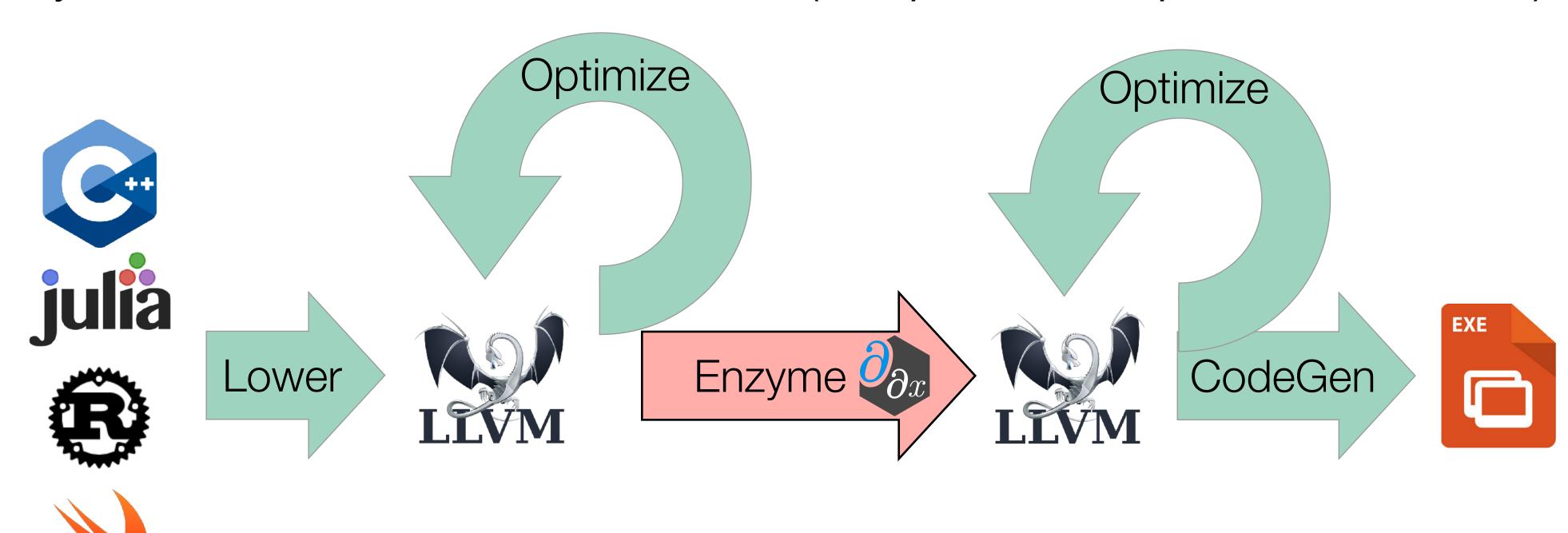
- Hinders application of ML to new domains
- Synthesizing gradients aims to close this gap





a Enzyme Overturns Conventional Wisdom

- As fast or faster than state-of-the-art tools
 - Running after optimization enables a 4.2x speedup
- Necessary semantics for AD derived at low-level (with potential cooperation of frontend)



Parallel Memory Detection

- Thread-local memory
 - Non-atomic load/store
- Same memory location across all threads
 - Parallel Reduction
- Others [always legal fallback]
 - Atomic increment

```
%tmp = load %d_res
store %d_res = 0
atomic %d_ptr += %tmp
```



Differentiation of SyncThreads

Case 3 [write sync write]

```
codeA(); // store %ptr
sync_threads;
codeB(); // store %ptr
diffe_codeB(); // load %d_ptr
               // store %d_ptr = 0
sync_threads;
diffe_codeA(); // load %d_ptr
               // store %d_ptr = 0
```

All uses of stores to d_ptr in diffe_B will correctly complete prior to diffe_A

Case 4 [read sync read]

```
codeA(); // load %ptr
sync_threads;
codeB(); // load %ptr
diffe_codeB(); // atomicAdd %d_ptr
sync_threads;
diffe_codeA(); // atomicAdd %d_ptr
```

Original and differential sync unnecessary and legal to include



CUDA Performance Improvements

- Introduce optimizations to reduce the use of memory
 - Alias Analysis to determine legality of recomputing an instruction
 - More aggressive alias analysis properties of syncthreads
 - Don't cache unnecessary values
 - Move cache outside of loops when possible
 - Heap-to-stack [and to register]
 - Don't cache memory itself acting as a cache [such as shared memory]
 - PHI Node unwrapping



Custom Derivatives & Multisource

One can specify custom forward/reverse passes of functions by attaching metadata

```
__attribute__((enzyme("augment", augment_func)))
__attribute__((enzyme("gradient", gradient_func)))
double func(double n);
```

 Enzyme leverages LLVM's link-time optimization (LTO) & "fat libraries" to ensure that LLVM bitcode is available for all potential differentiated functions before AD

